# Stable nonlinear manifold approximation using compositional networks

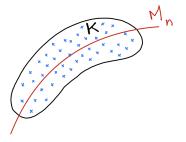
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Joint works with Antoine Bensalah and Joel Soffo.

#### Introduction

We consider the problem of approximating a subset K of a normed space X by a low-dimensional set  $M_n$ , from samples in K.



A common setting (in statistics) if when K is the range of some vector or function-valued random variable.

Another classical setting is the solution of forward or inverse problems for parameter-dependent equations, where

$$K = \{u(y) : y \in Y\}$$
 with  $R(u(y); y) = 0$ .

#### Introduction

#### The approximating sets $M_n$ can be

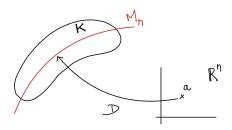
- $\bullet$  constructed offline by manifold approximation methods, using samples from K,
- used online to compute approximations of elements in K with low computational complexity, or from limited information.

#### **Encoder-Decoder**

A large class of manifold approximation methods can be described by an encoder  $E: K \to \mathbb{R}^n$  and a decoder  $D: \mathbb{R}^n \to X$ .

The decoder provides a parametrization of a *n*-dimensional "manifold"

$$M_n = \{D(a) : a \in \mathbb{R}^n\}.$$



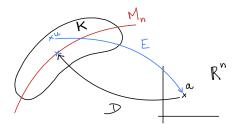
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The encoder is related to the approximation process (algorithm). It associates to  $u \in K$  a parameter value  $a = E(u) \in \mathbb{R}^n$ .



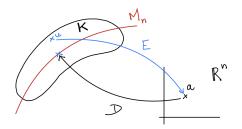
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An element  $u \in K$  is approximated by  $D \circ E(u) \in M_n$ .

This problem is equivalent to approximating the identity map on K by  $D \circ E$  (auto-encoder of K).

## **Optimal performance**

Manifold approximation methods can be classified in terms of the properties of their encoders and decoders.

The optimal performance of a given class  $\mathcal{E}_n$  of encoders from X to  $\mathbb{R}^n$  and a given class  $\mathcal{D}_n$  of decoders from  $\mathbb{R}^n \to X$  can be assessed in worst-case setting by

$$\inf_{D\in\mathcal{D}_n,E\in\mathcal{E}_n}\sup_{u\in K}\|u-D\circ E(u)\|_X.$$

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If the set K is equipped with a measure  $\rho$ , the optimal performance can be measured in average sense by

$$\inf_{D\in\mathcal{D}_n,E\in\mathcal{E}_n} \left( \int_K \|u-D\circ E(u)\|_X^p d\rho(u) \right)^{1/p}.$$

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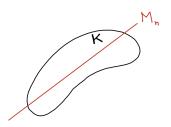
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These errors define measures of complexity (widths) of K.

The range  $M_n$  of a linear decoder  $D: \mathbb{R}^n \to X$  is a linear space with dimension at most n.

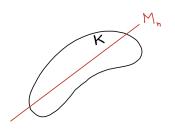


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Restricting the decoder and encoder to be linear yields the approximation numbers (linear widths)

$$a_n(K)_X = \inf_{rank(A)=n} \sup_{u \in K} \|u - Au\|_X,$$

where the infimum is taken over all linear maps  $A: X \to X$  with rank n.

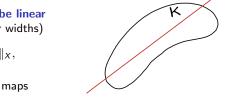


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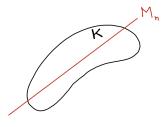
$$d_n(K)_X = \inf_{D \in L(\mathbb{R}^n;X)} \sup_{u \in K} \inf_{a \in \mathbb{R}^n} \|u - D(a)\|_X = \inf_{\dim M_n = n} \sup_{u \in K} \inf_{v \in M_n} \|u - v\|_X.$$

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For X a Hilbert space,  $a_n(K)_X = d_n(K)_X$  and an optimal auto-encoder  $D \circ E$  is given by the orthogonal projection  $P_{M_n}$  onto an optimal space  $M_n$ .

In practice optimal linear spaces in worst-case error are out of reach but near to optimal spaces  $M_n$  can be obtained by greedy algorithms, that generate an increasing sequence of spaces from samples in K [Buffa, Maday, Patera, Prud'homme, Turinici, Binev, Cohen, Dahmen, DeVore,...].

## Linear approximation - Average setting

When X is a Hilbert space and the error is measured in average sense, it yields the average Kolmogorov n-width

$$d_n^{(p)}(K,\rho)_X^p = \inf_{D \in \mathcal{L}(\mathbb{R}^n;X)} \int_K \inf_{a \in \mathbb{R}^n} \|u - D(a)\|_X^p d\rho(u) = \inf_{\dim M_n = n} \int_K \|u - P_{M_n}u\|_X^p d\rho(u).$$

For p=2 (mean-squared error), an optimal space  $M_n$  is given by a dominant eigenspace of the (compact) operator

$$T(v) = \int_K u(u,v)_X d\rho(u), \quad v \in X,$$

and

$$\int_{K} \|u - P_{\mathbf{M_n}} u\|_{X}^{2} d\rho(u) = d_n^{(2)}(K, \rho)_{X}^{2} = \sum_{i > n} \lambda_i(T)$$

If  $\rho$  is a probability measure, assuming  $\bar{u} = \int u \, d\rho(u) = 0$ , T is the covariance operator of  $\rho$  and this corresponds to Principal Component Analysis (PCA).

## Nonlinear continuous manifold approximation

Restricting both the encoder and decoders to be continuous (possibly nonlinear) yields the notion of nonlinear manifold width of [DeVore, Howard and Michelli 1989]

$$\delta_n(K)_X = \inf_{E \in C(X;\mathbb{R}^n)} \inf_{D \in C(\mathbb{R}^n;X)} \sup_{u \in K} \|u - D \circ E(u)\|_X.$$

Further restricting encoders and decoders to be Lipschitz continuous yields the notion of stable manifold width [Cohen et al 2022]

$$\delta_n^L(K)_X = \inf_{\text{Lip}(E) \le L} \inf_{\text{Lip}(D) \le L} \sup_{u \in K} \|u - D \circ E(u)\|_X.$$

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Lipschitz continuity ensures stability of the approximation process, a crucial property in practice.

However, implementation of optimal nonlinear encoders may be difficult or even infeasible, e.g. associated with NP-hard optimization problems.

Restricting the encoder to be linear and continuous yields the *n*-th minimal error of linear information (sometimes called sensing numbers)

$$e_n(K)_X = \inf_{\substack{D \ \ell_1, \dots, \ell_n \ u \in K}} \sup_{\substack{u \in K}} \|u - D(\ell_1(u), \dots, \ell_n(u))\|_X$$

where the infimum is taken over all linear forms  $\ell_1, \ldots, \ell_n$  and all nonlinear maps D.

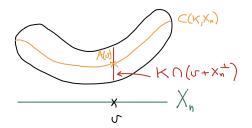
This benchmark are relevant in many applications where the available information  $E(u)=(\ell_1(u),\ldots,\ell_n(u))$  is linear in u (point evaluations of functions, local averages of functions or more general linear functionals).

In a Hilbert setting and worst case setting, this corresponds to

$$\inf_{\dim(X_n)=n}\inf_{A:X_n\to X}\sup_{u\in K}\|u-A(P_{X_n}u)\|_X$$

For a given  $X_n$ , an optimal algorithm  $A: X_n \to X$  is such that A(v) is the Chebychev center of the slice

$$K \cap (v + X_n^{\perp}) = \{u \in K : P_{X_n} u = v\}.$$



This yields an optimal n-dimensional manifold

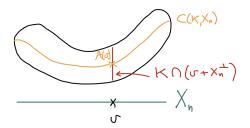
$$M_n = C(K, X_n) := \{cen(K \cap (v + X_n^{\perp})) : v \in P_{X_n}K\}.$$

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In average setting, an average minimal error can be defined, and optimal manifold is obtained by averaging over slices.

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With  $(\varphi_i)_{i\geq 1}$  an orthonormal basis of X and  $X_n=span\{\varphi_1,\ldots,\varphi_n\}$ , this is associated with the linear encoder

$$E(u) = ((u, \varphi_i))_{i=1}^n$$

and a decoder of the form

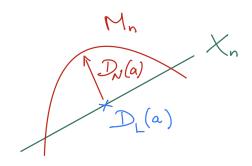
$$D(a) = A(\sum_{i=1}^{n} a_i \varphi_i) = \sum_{i=1}^{n} a_i \varphi_i + \sum_{i>n} g_i(a) \varphi_i = D_L(a) + D_N(a),$$

where the functions  $g_i : \mathbb{R}^n \to \mathbb{R}$  are nonlinear maps.

 $D_L$  is a linear operator from  $\mathbb{R}^n$  to  $X_n := span\{\varphi_1, \dots, \varphi_n\}$  such that  $D_L(E(u)) = P_{X_n}u$ .

 $D_N$  maps  $\mathbb{R}^n$  to the complementary space of  $X_n$  in X.

For a feasible implementation, we truncate to the first N terms, so that the range of D is a nonlinear manifold  $M_n \subset X_N = span\{\varphi_1, \ldots, \varphi_N\}.$ 



The structure of the decoder relies on the fact that for

$$u = \sum_{i=1}^{n} a_i(u)\varphi_i + \sum_{i>n} a_i(u)\varphi_i \in K,$$

the coefficients  $a_i(u)$  for i > n may be well approximated as functions  $g_i(E(u))$  of the first few coefficients  $E(u) = (a_i(u))_{i=1}^n$ .

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Natural choices for functions  $g_i$  are

- Quadratic polynomials [Barnett and Farhat 2022][Geelen et al 2023]
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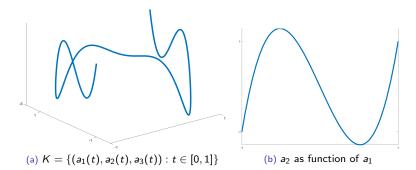
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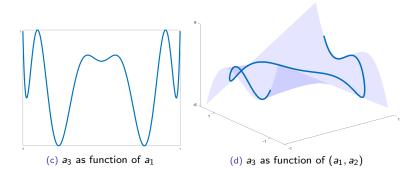
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The relation between  $a_i(u)$  and E(u) may be highly nonlinear. Even highly expressive approximation tools may result in poor accuracy, due to the difficulty of learning with limited data.

In many applications, a coefficient  $a_i(u)$  for i > n may have a highly nonlinear relation with the first n coefficients a = E(u) but a much smoother relation when expressed in terms of a and additional coefficients  $a_j(u)$  with n < j < i.



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# Decoder based on compositional polynomial network (CPN)

This suggests the following compositional structure of the decoder's functions

$$g_i(a) = f_i(a, (g_j(a))_{n < j \le n_i}),$$

where the  $f_i$  are polynomial functions.

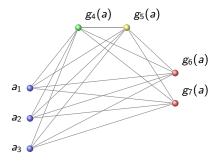


Figure: A compositional polynomial network (CPN) with N=7 and n=3, maximum number of compositions 3.

# Decoder based on compositional polynomial network (CPN)

The variables  $(a, (g_j(a))_{n < j \le n_i})$  take values in a set of measure zero in  $\mathbb{R}^{n_i}$ , but it is still possible to learn polynomial functions  $f_i$  from a limited training set.

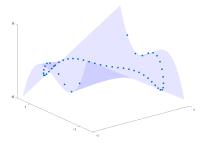


Figure: Learning  $a_3$  as function of  $(a_1, a_2)$ 

In practice, for high-dimensional approximation of  $f_i: \mathbb{R}^{n_i} \to \mathbb{R}$  from samples, use of sparse polynomial approximation (CPN-S) or low-rank approximation (tensor networks) (CPN-LR) in  $\mathbb{P}_p^{\otimes n_i}$ .

# Control of error (mean-squared setting)

Assume X is a Hilbert space and consider  $D: \mathbb{R}^n \to X_N$ , with  $X_N$  the subspace with orthonormal basis  $\varphi_1, \ldots, \varphi_N$ .

The mean squared error

$$e_2(D \circ E)^2 := \|id - D \circ E\|_2^2 := \int_{\kappa} \|u - D(E(u))\|_X^2 d\rho(u)$$

satisfies 
$$e_2(D \circ E)^2 = \sum_{i=n+1}^N \epsilon_{i,2}^2 + \|id - P_{X_N}\|_2^2, \quad \epsilon_{i,2} = \|a_i - g_i \circ a\|_2$$

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Given a prescribed precision  $0 < \epsilon < 1$ , it holds

$$e_2(D \circ E) \leq \epsilon e_2(0)$$

whenever

$$\|id - P_{X_N}\|_2^2 \le \beta \epsilon^2 e_2(0)^2$$
 (1)

$$\epsilon_{i,2}^2 \le \overline{\epsilon}_{i,2}^2 := \omega_i (1 - \beta) \epsilon^2 e_2(0)^2, \quad \forall i > n$$
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- (1) is achieved by using PCA to define  $X_N$ , with a suitable selection of  $N = N(\epsilon)$ .
- (2) is achieved by a control of the approximation of  $a_i$  by  $f_i((g_j(a))_{j=1}^{n_i})$ , using validation.

## Control of stability

Given an orthonormal basis  $\varphi_1, \ldots, \varphi_N$ , the encoder  $E: X \to \mathbb{R}^n$  is 1-Lipschitz

$$\boxed{\|E(u) - E(u')\|_2 = \|P_{X_n}(u - u')\|_X \le \|u - u'\|_X}$$

Given  $\gamma = (\gamma_i)_{i=n+1}^N$ , we equip  $\mathbb{R}^{n_i}$  with the norm

$$||b||_{i,\gamma} = \max\{||(b_j)_{j=1}^n||_2, \max_{n < j \le n_i} \gamma_j^{-1}|b_j|\}$$

and define the corresponding Lipschitz norm (estimated from samples)

$$||f_i||_{i,\gamma} = \max_{b,b'} \frac{|f_i(b) - f_i(b')|}{||b - b'||_{i,\gamma}}$$

Given a prescribed Lipschitz constant  $L \ge 1$ , letting  $||f_i||_{i,\gamma} = \gamma_i$  and assuming  $\gamma_i^2 \le \overline{\gamma}_i^2$  with  $\sum_{i=n+1}^n \overline{\gamma}_i^2 \le L^2 - 1$ , it holds

$$||D(a) - D(a')||_X \le L||a - a'||_2$$

## Adaptive algorithm

The prescribed bounds for errors

$$\epsilon_{i,p} \leq \overline{\epsilon}_{i,p}$$

or Lipschitz constants

$$\gamma_i = ||f_i||_{i,\gamma} \leq \bar{\gamma}_i$$

may not be satisfied for some indices  $i \in \{1, ..., N\}$ .

This requires to progressively adapt the set of indices  $\{1, \ldots, n\}$  associated with the encoder.

Prescribed upper bounds  $\bar{\epsilon}_{i,p}$  and  $\bar{\gamma}_i$  can be updated during the algorithm in order to obtain a sharper control of error and stability.

#### Numerical illustration: KdV

We consider the Korteweg-de Vries (KdV) equation

$$\frac{\partial u}{\partial t} + 4u \frac{\partial u}{\partial x} + \frac{\partial^3 u}{\partial x^3} = 0 \quad on \quad [-\pi, \pi] \times [0, 1]$$

with periodic boundary conditions and some initial condition. We consider the manifold

$$K = \{u(\cdot, t) : t \in [0, 1]\}$$

We use 5000 samples  $u_i = u(\cdot, t_i)$  as training samples.

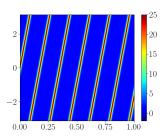


Figure: Function u(x, t).

Method	р	n	N	REtrain	RE <sub>test</sub>
Linear	/	5	5	0.435	0.450
Quadratic	2	5	20	0.094	0.099
Additive + AM	5	5	43	0.081	0.084
Sparse	5	5	43	0.013	0.014
CPN-S ( $\varepsilon = 10^{-4}$ )	5	5	43	0.000072	0.000074

Table: Comparison of methods for the same manifold dimension n = 5. CPN-S with sparse polynomials of degree p = 5.

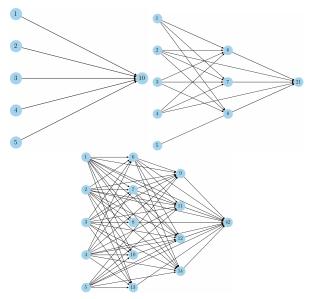


Figure: Compositional networks for  $g_{10},\,g_{21}$  and  $g_{42}$  ( $\epsilon=10^{-4},\,p=5$ )

Tolerance	n	Ν	RE <sub>train</sub>	RE <sub>test</sub>
$\varepsilon = 10^{-1}$	2	15	$6.2 \times 10^{-2}$	$6.4 \times 10^{-2}$
$\varepsilon = 10^{-2}$	2	25	$6.67 \times 10^{-3}$	$6.84 \times 10^{-3}$
$\varepsilon = 10^{-3}$	3	34	$6.83 \times 10^{-4}$	$7 \times 10^{-4}$
$\varepsilon=10^{-4}$	5	43	$7.170 \times 10^{-5}$	$7.367 \times 10^{-5}$
$\varepsilon = 10^{-5}$	6	52	$6.76  imes 10^{-6}$	$6.916 \times 10^{-6}$
$\varepsilon=10^{-6}$	11	61	$7.689 \times 10^{-7}$	$7.885 \times 10^{-7}$

Table: Results of CPN-S for p=5 and different target precisions  $\varepsilon$ .

р	n	Ν	RE <sub>train</sub>	$RE_{test}$
3	9	43	$7.401 \times 10^{-5}$	$7.574 \times 10^{-5}$
4	7	43	$7.524 \times 10^{-5}$	$7.720 \times 10^{-5}$
5	5	43	$7.170 \times 10^{-5}$	$7.367 \times 10^{-5}$

Table: Results of CPN-S with different degrees p for  $\epsilon=10^{-4}$ 

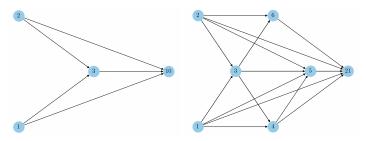


Figure: Compositional networks for  $g_{10}$  and  $g_{21}$  ( $\epsilon=10^{-2}$ , p=5)

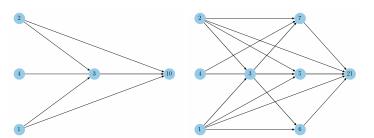


Figure: Compositional networks for  $g_{10}$  and  $g_{21}$  ( $\epsilon=10^{-3},\ p=5$ )

Method	р	n	Ν	$RE_{train}$	RE <sub>test</sub>
Linear	/	2	/	$6.63 \times 10^{-1}$	$6.91 \times 10^{-1}$
Quadratic	2	2	5	$5.33 \times 10^{-1}$	$5.66 \times 10^{-1}$
Additive-AM	5	2	43	$2.10 \times 10^{-1}$	$5.47 \times 10^{-1}$
Sparse	5	2	43	$1.73 \times 10^{-1}$	$1.85\times10^{-1}$
Low-Rank	5	2	43	$7.47 \times 10^{-2}$	$7.94 \times 10^{-2}$
CPN-LR $(\epsilon=10^{-4})$	5	2	43	$6.85 \times 10^{-5}$	$7.06 \times 10^{-5}$

Table: Comparison of methods for the same manifold dimension n=5. CPN-LR uses low-rank polynomials with degree p=5 (tensor train format).

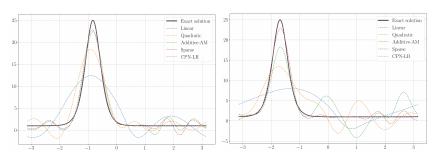


Figure: Comparison of methods for predicting  $u(\cdot,t)$  at t=0.5 (left) and t=1 (right). Same dimension n=5.

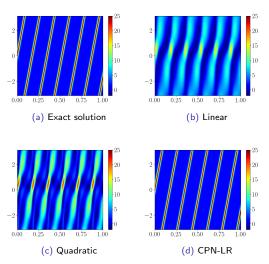


Figure: Predictions for different methods, with manifold dimension n = 2.

## Optimal linear encoders and associated spaces

The encoder (or associated space  $X_n$ ) can be optimized.

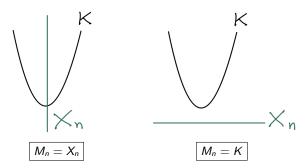
A natural choice is to consider the optimal or near-optimal spaces of linear methods, given by principal component analysis (optimal in mean-squared error) or greedy algorithms (close to optimal in worst case error) [Barnett and Farhat 2022, Geelen et al 2023, Barnett, Farhat and Maday 2023, Geelen et al 2024].

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However, these may be far from optimal for nonlinear approximation.

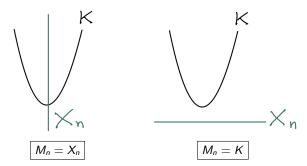


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In [Schwerdtner and Peherstorfer 2024], greedy algorithm for a (near to) optimal construction of  $X_n$  for quadratic manifold approximation (quadratic decoder).

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The benchmark for a linear encoder and a Lipschitz stable decoder is

$$e_n^L(K)_X = \inf_{Lip(E) \le 1} \inf_{Lip(D) \le L} \sup_{u \in K} ||u - D(E(u))||_X$$

where the infimum is taken over linear 1-Lipschitz encoders and L-Lipschitz decoders.

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$$e_n^L(K)_X = \inf_{\dim(X_n) = n} \inf_{Lip(A) < L} \sup_{u \in K} \|u - A(P_{X_n}u)\|_X$$

where

$$Lip(A) = \sup_{\substack{u,v \in K \\ P_{X_n}(u-v) \neq 0}} \frac{\|A(P_{X_n}u) - A(P_{X_n}u)\|_X}{\|P_{X_n}u - P_{X_n}v\|_X} \leq L$$

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A dual version is

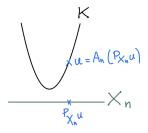
$$\tilde{e}_n^{\eta}(K)_X = \inf_{\dim(X_n)=n} \inf_A Lip(A)$$

where the infimum is taken over maps A such that

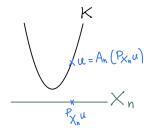
$$\sup_{u\in K}\|u-A(P_{X_n}u)\|_X\leq \eta$$

 $e_n^L(K)_X$  or  $\tilde{e}_n^{\eta}(K)_X$  can be seen as notions of "Lipschitz widths" of K, using linear encoder. See related notion in [Petrova and Wojtaszczyk 2023] with arbitrary encoders.

Consider  $\eta = 0$ , assuming the existence of an exact recovery map  $A_n$  for some  $X_n$ . This requires K to be a n-dimensional manifold, and  $A_n$  is a global chart associated to  $X_n$ .



Consider  $\eta = 0$ , assuming the existence of an exact recovery map  $A_n$  for some  $X_n$ . This requires K to be a n-dimensional manifold, and  $A_n$  is a global chart associated to  $X_n$ .



Then we obtain a measure of Lipschitz regularity of K related to  $X_n$  as

$$Lip(A_n) = \sup_{\substack{u,v \in K \\ P_{X_n}(u-v) \neq 0}} \frac{\|u-v\|_X}{\|P_{X_n}(u-v)\|_X} =: Lip(K, X_n)_X$$

and the corresponding notion of Lipschitz width of K

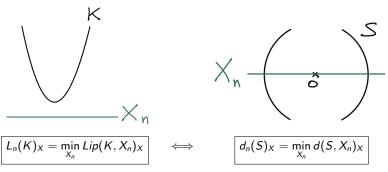
$$L_n(K)_X := \inf_{\dim(X_n)=n} Lip(K, X_n)_X$$

which defines an optimal space  $X_n$  in terms of stability of the corresponding chart  $A_n$ .

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It holds

$$L_n(K)_X = \frac{1}{1-d_n(S)_X} \quad \text{with} \quad S = \{\frac{z}{\|z\|_X} : z \in (K-K) \setminus \{0\}\}.$$



Close to optimal spaces can be obtained by greedy algorithms.

In the average setting, possible relaxation by considering the average width  $d^{(2)}(S,\pi)_X$ , with  $\pi$  the push-forward measure on S of  $\rho\otimes\rho$  through the map  $(u,v)\mapsto (u-v)/\|u-v\|_X$ .

Instead of

$$d_n(S) = \inf_{\dim(X_n) = n} \sup_{z \in S} ||z - P_{X_n}z||_X$$

we consider

$$d^{(2)}(S,\pi)_X = \inf_{\dim(X_n)=n} \int_S \|z - P_{X_n}z\|_X^2 d\pi(z)$$

that is an eigenvalue problem.

In practice, we introduce estimators from finitely many samples in K.

## Optimal spaces guided by Lipschitz stability

Previous notions may not be well defined for sets K that are not smooth n-dimensional manifolds described by global chart.

This requires to consider relaxations  $\tilde{e}_n^n(K)_X$  that balance Lipschitz stability and approximation error, or *n*-dimensional approximations of the set K.

The design of practical strategies using samples from K is a challenging problem.

### **Conclusions**

 Nonlinear manifold approximation method with linear encoder and nonlinear decoder based on compositional polynomial networks, with control of error and stability [1].

#### Reference

[1] A. Bensalah, A. Nouy, and J. Soffo. Nonlinear manifold approximation using compositional polynomial networks. arXiv e-prints arXiv:2502.05088, Feb. 2025.

### **Conclusions**

- Nonlinear manifold approximation method with linear encoder and nonlinear decoder based on compositional polynomial networks, with control of error and stability [1].
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#### **Conclusions**

- Nonlinear manifold approximation method with linear encoder and nonlinear decoder based on compositional polynomial networks, with control of error and stability [1].
- Optimal spaces/encoders based on Lipschitz stability of decoders.
- Some challenging problems related to online approximation with nonlinear manifolds: optimization and sampling/discretization.

#### Reference

[1] A. Bensalah, A. Nouy, and J. Soffo. Nonlinear manifold approximation using compositional polynomial networks. arXiv e-prints arXiv:2502.05088, Feb. 2025.

## THANK YOU FOR YOUR ATTENTION

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